

Lecture 38

About Discrete Math, 327, Fall.

~~∇~~ too much overlap with
Math 318

Finite Algebraic Systems will
be studied.

page,

- Degree of a vertex, regular graph, even and odd vertices.
- Degree sequences. The handshake lemma.
- Isomorphism.
- Classifying small graphs.
- Adjacency matrices
- Incidence matrices.
- Walks, Trails, Paths.
- Bipartite graphs.
- Cube graphs. The Petersen graph. Trees. Cycles.
- Rooted trees.
- Directed graphs
- Orientability. Bridges.
- Adjacency and Incidence matrices for digraphs
- Counting digraphs
- In-degree and out-degree sequences.
- Finite State Automata and their languages.
- Euler circuits.
- Edge traceable graphs
- Weighted digraphs.
- Minimal cost paths.
- Longest paths.
- Degrees of connectedness.
- Cut-sets and vertex-cut-sets.
- Menger's Theorem
- Planar graphs. Euler's theorem.
- Coloraries to Euler's Thereom
(have degree ≤ 5 vertex)
- ($m \leq 3n - 6$) & tori
- Graphs drawn on Klein bottles.

Final will evenly cover all parts of the class.

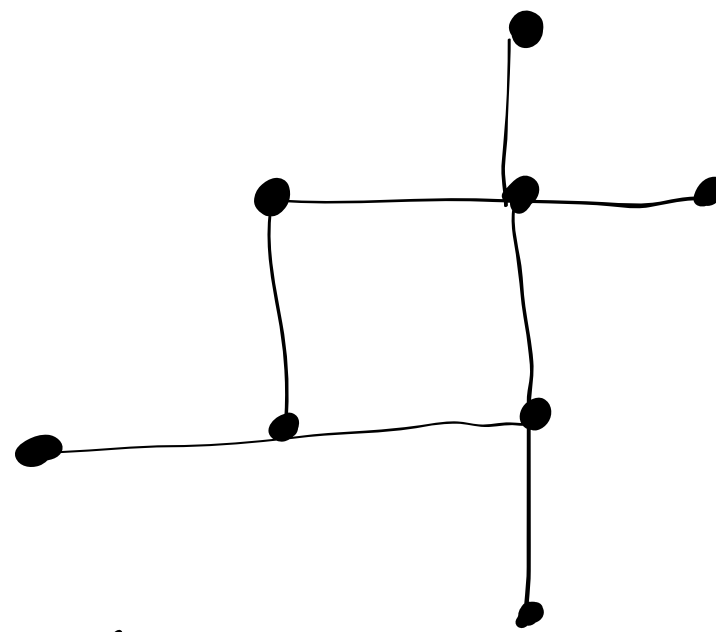
No applications except FSA

Final. The final exam is Monday, May 6, 10:00-12:00. Web student's not attending this exam need to set a time with me.

Dual Graphs

page 2

#1 Consider G :



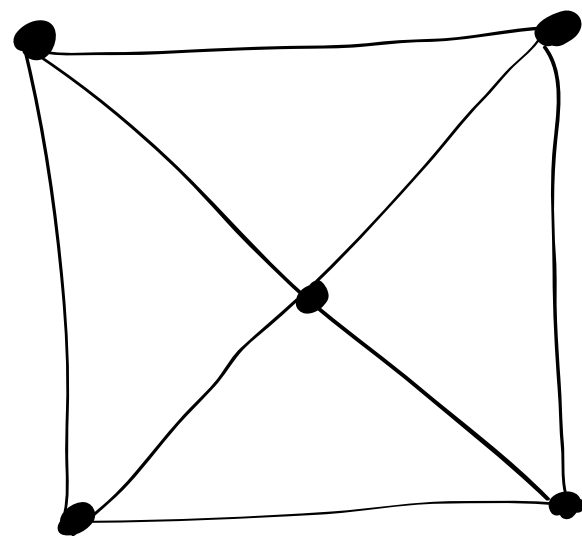
Some sample problems.

(a) What is the chromatic number of G ?

(b) How many ways can G be colored with k -colors, when k is the chromatic number of G ?

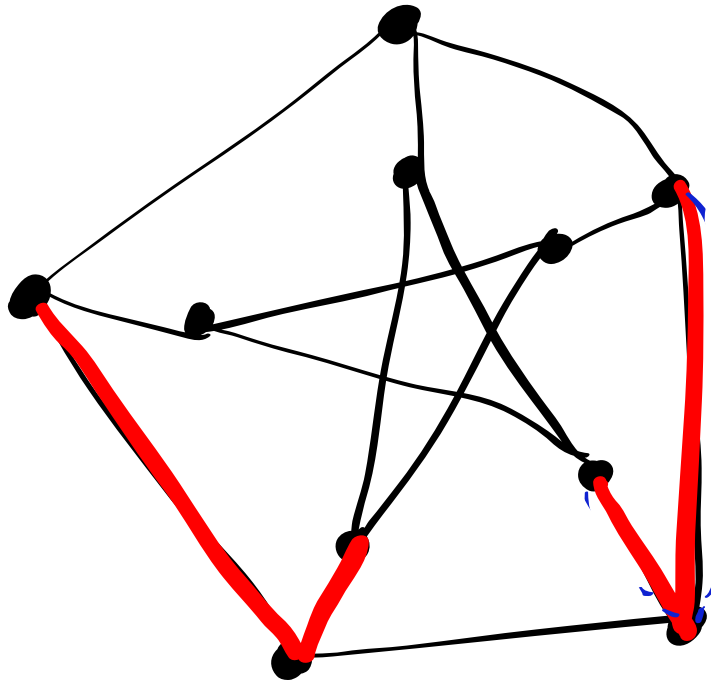
#2.

Let H be
Is the dual of
 H edge-traceable?



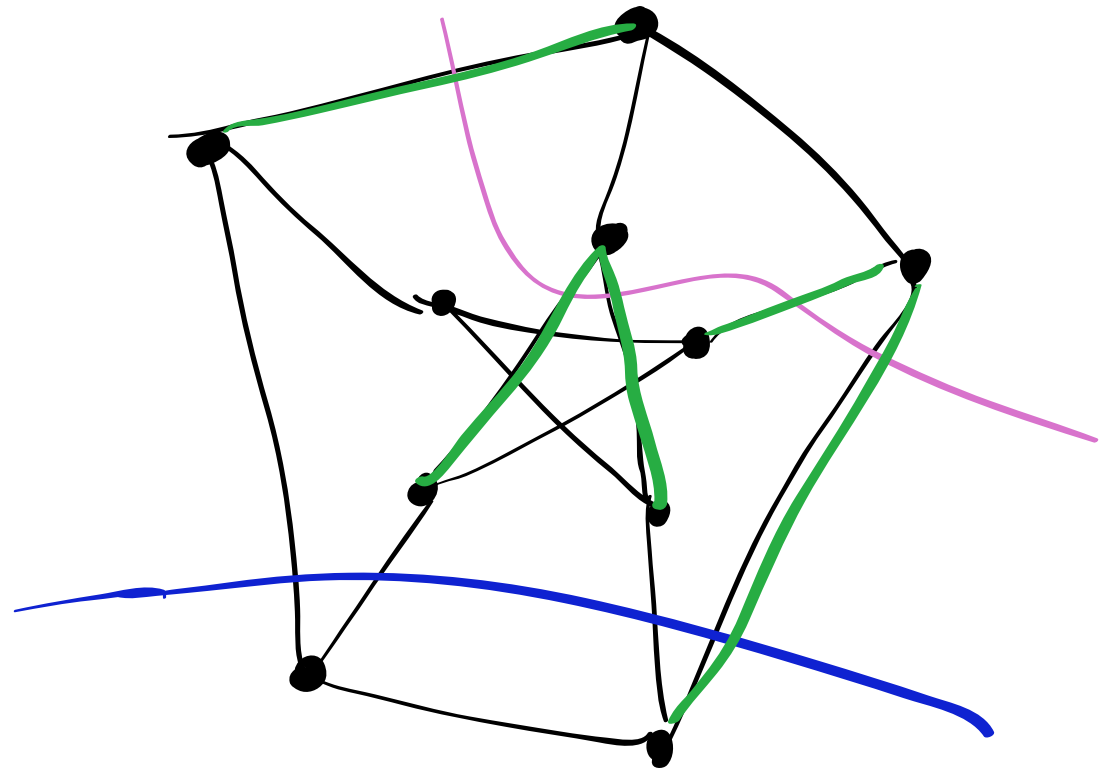
page 3

edge-cut-sets

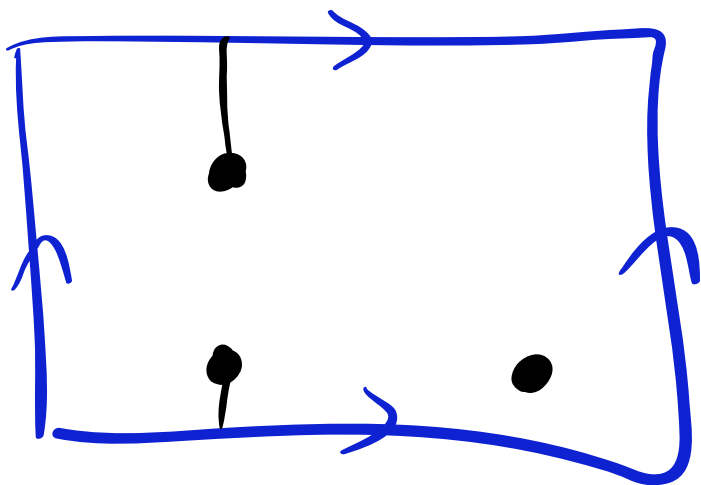
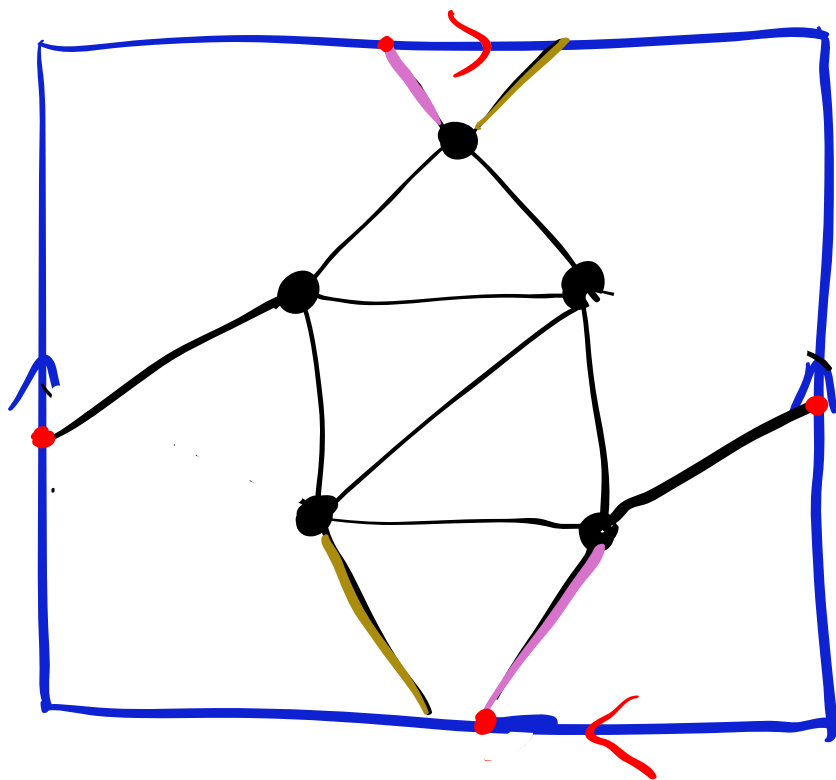


Find an edge-cut-set
with ~~4~~ vertices.

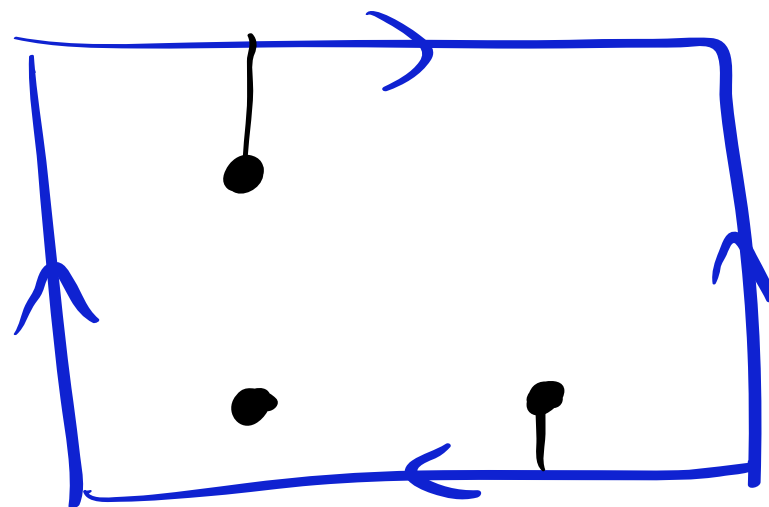
5



Draw K_5 on a Torus.



Torus



Klein Bottle

Show no simple graph G of 11 vertices
can be planar if \bar{G} is planar.

$$G. \quad \begin{array}{l} n=11 \\ m=? \end{array}$$

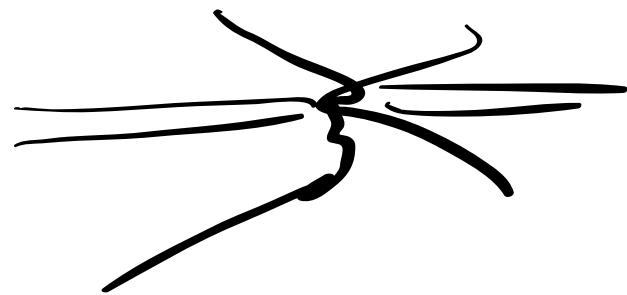
$$\bar{G} \quad \begin{array}{l} n=11 \\ m' = 55 - m \end{array}$$

$$\frac{11 \cdot 10}{2} = 55$$

Cor. to Euler: $\left\{ \begin{array}{l} m \leq 3 \cdot 11 - 6 = 27 \\ m' \leq 3 \cdot 11 - 6 = 27 \end{array} \right.$

If both are planar \rightarrow
+

$$55 = m + m' \leq 54$$

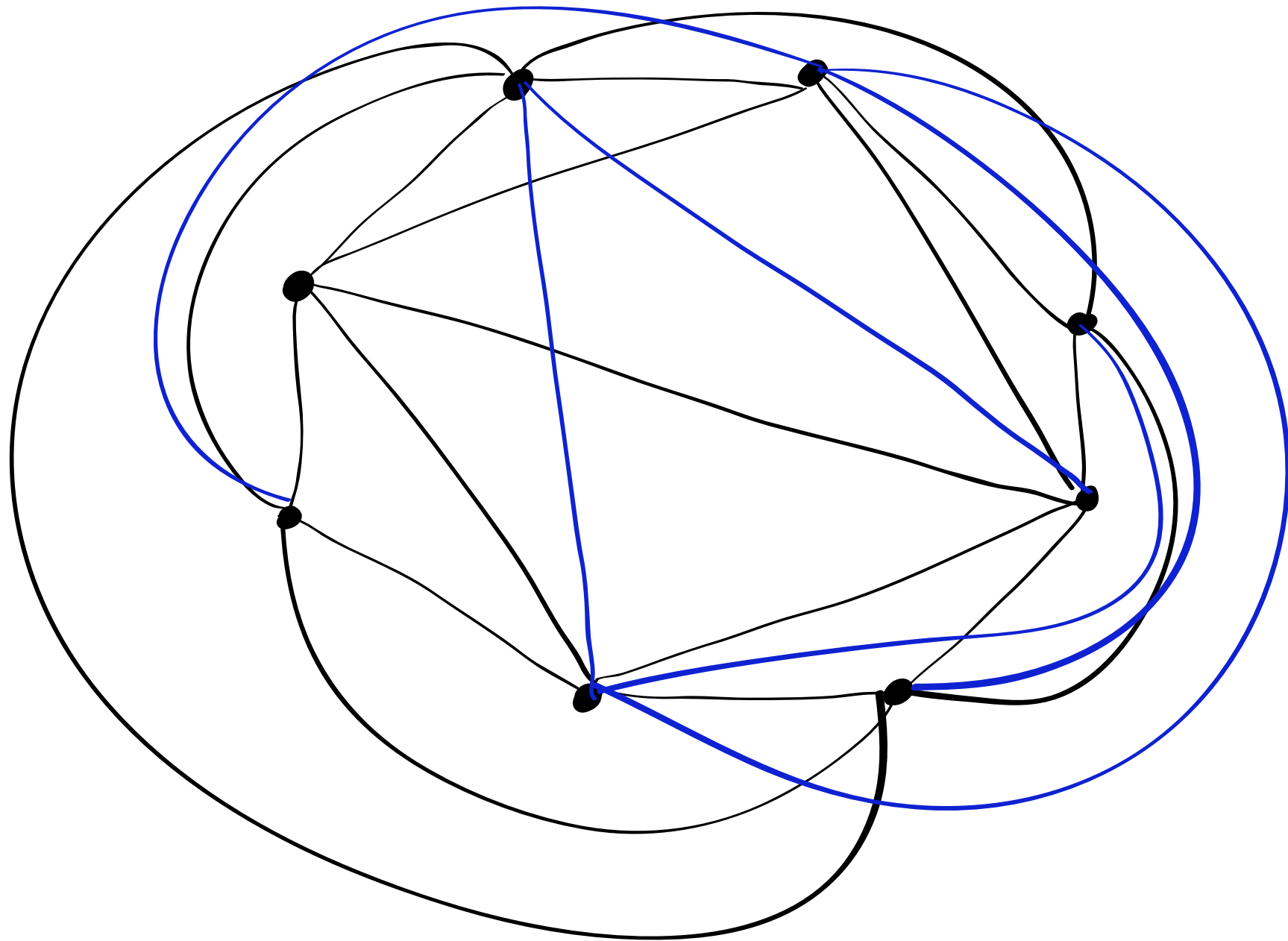


$n = 8$: G, \bar{G} planar; $m + m' = \frac{8 \cdot 7}{2} = 28$

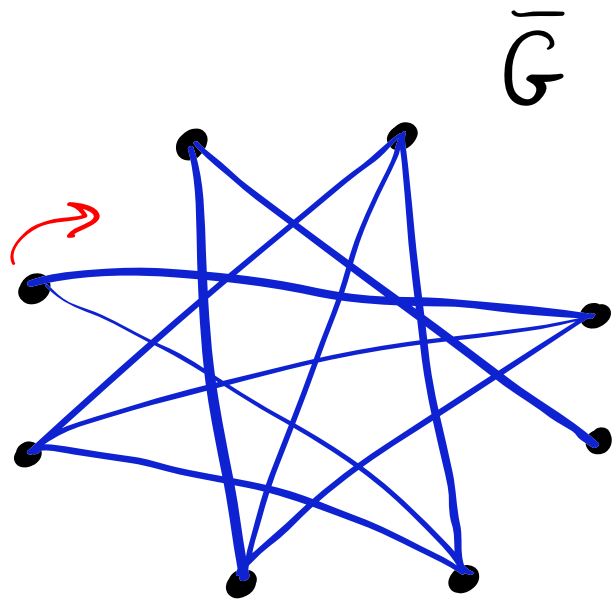
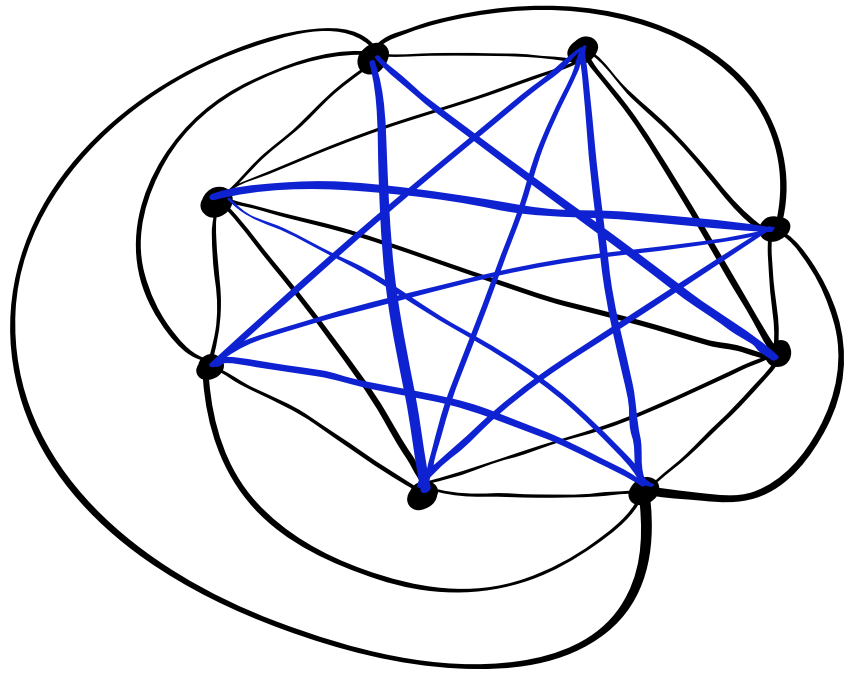
$m \leq 3 \cdot 8 - 6 = 18$

$m' \leq 18$

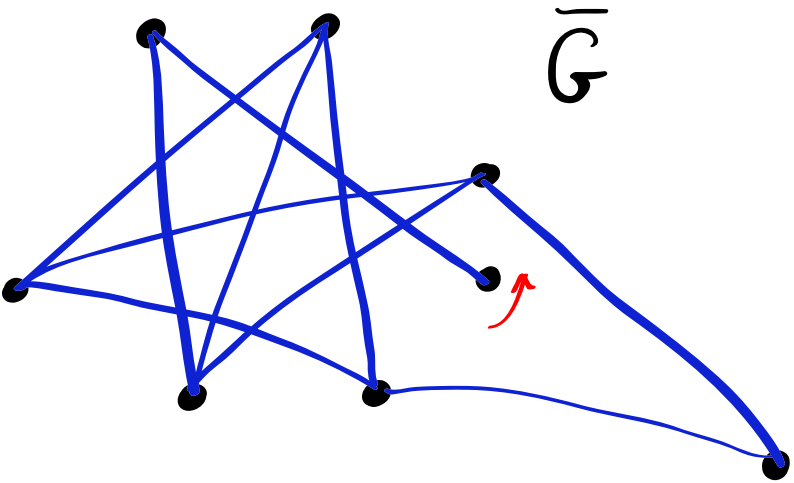
G with ≈ 18 edges:



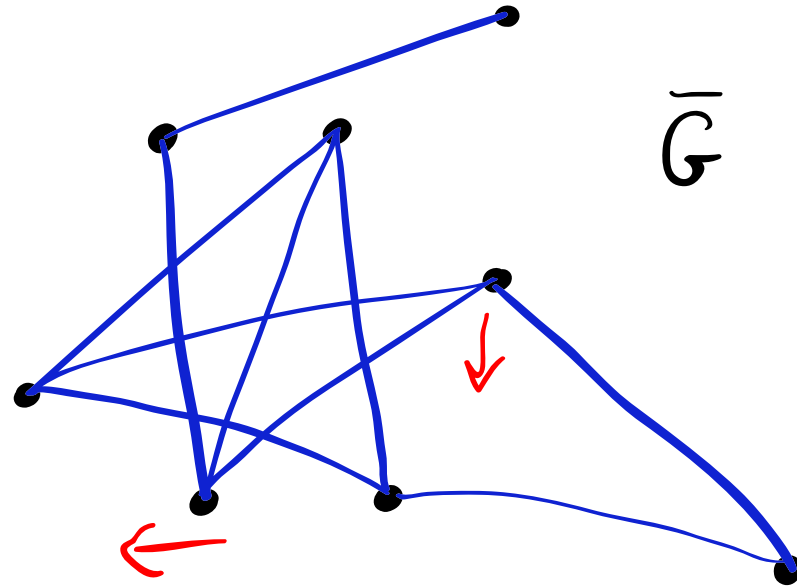
~~etc.~~
dang!



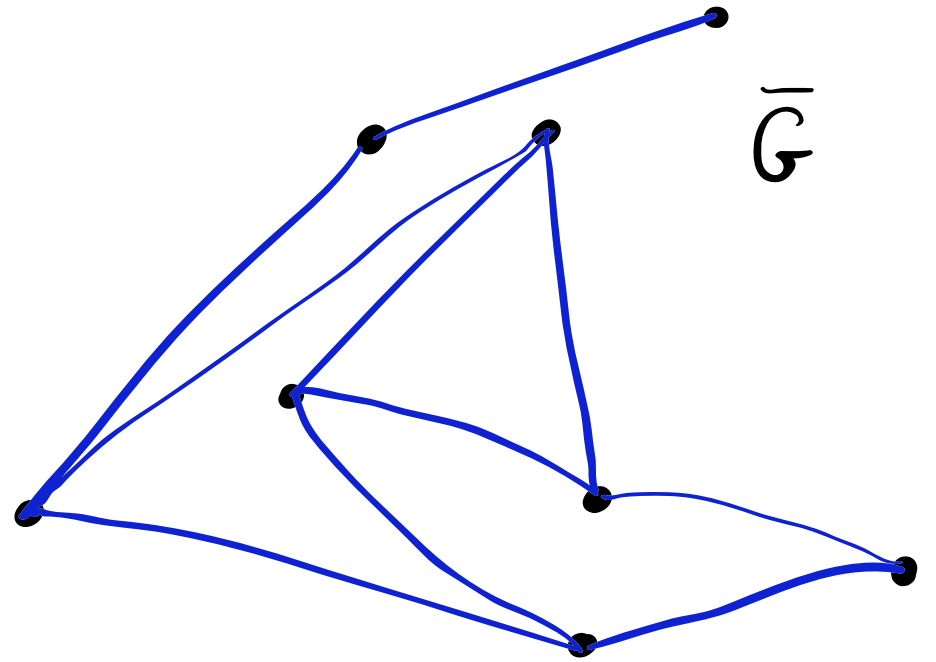
G_1



G_1



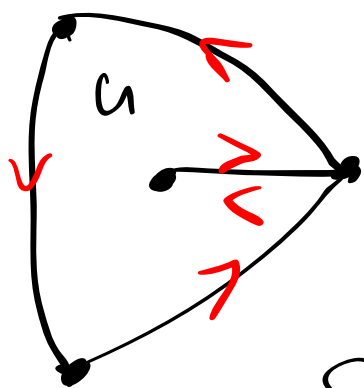
G_1



G_1

Girth = length of shortest cycle
 in G . if G has girth g and is
 planar (simple, connected) then $m \leq \frac{(n-2)g}{g-2}$.

The boundary of a face is closed path.



Therefore, if u , a face, has
 degree d , then $d \geq g$.

$$\text{So } 2m = \sum_{u \text{ a face}} \text{deg}(u) \geq \sum g = fg.$$

$$\therefore f \leq \frac{2m}{g}. \quad 2 = n - m + f \leq n - m + \frac{2m}{g}.$$

$$\Rightarrow m - \frac{2m}{g} \leq n - 2 \Rightarrow m \leq g \frac{n-2}{g-2}.$$